The Bossa Framework for Scheduler Development

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Target domain: Kernel process scheduling

Process scheduling: How an OS selects a process for the CPU.

- Many scheduling policies (round-robin, RM, EDF, etc.).
- No policy is perfect for all applications.

Implementing a scheduler requires:

- Understanding the scheduling policy.
- Understanding the target OS.
 - Any error can crash the machine.

Our proposal: Bossa

A Run-Time System for integrating scheduling policies into a legacy OS

- Generates event notifications.
- Defines a model of kernel scheduling requirements.

A Domain-Specific Language for implementing scheduling policies.

- DSL: A language dedicated to a particular domain.
- Provides high-level domain-specific abstractions that
 - Hide technical details.
 - Simplify programming.
 - Enable verifications, optimizations.



- The process scheduling problem.
- The Bossa DSL.
- Preparing Linux for use with Bossa.
- Verification.
- Performance.
- Conclusion, ongoing work.

CPU:

Other processes:

CPU:

Other processes:

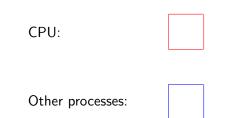
A process arrives.

CPU:



Other processes:

The process is elected.



Another process arrives.

CPU:

Other processes:



The red process blocks.

CPU:



Other processes:



The blue process is elected.

CPU:



Other processes:



Another process arrives.

CPU:



Other processes:



The red process unblocks.

CPU:

Other processes:



The blue process blocks.

CPU:

Other processes:



Which process is elected next?

- Process states (running, ready, blocked, etc.).
- Election criteria.
- OS events (blocking, unblocking, etc.).

```
scheduler EDF = {
```

states = {
 running
 ready
 blocked
 computation_ended
 terminated

}

}

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```
scheduler EDF = {
```

```
states = {
  RUNNING running
  READY ready
  BLOCKED blocked
  BLOCKED computation_ended
  TERMINATED terminated
}
```

}

```
scheduler EDF = {
```

```
states = {
  RUNNING running : process;
  READY ready : select queue;
  BLOCKED blocked : queue;
  BLOCKED computation_ended : queue;
  TERMINATED terminated;
}
```

```
scheduler EDF = {
```

```
states = {
  RUNNING running : process;
  READY ready : select queue;
  BLOCKED blocked : queue;
  BLOCKED computation_ended : queue;
  TERMINATED terminated;
}
ordering_criteria = {
```

}

```
scheduler EDF = {
  process = {
    time period;
   time wcet;
    time current_deadline;
   timer period_timer;
  }
  states = {
    RUNNING running : process;
    READY ready : select queue;
    BLOCKED blocked : queue;
    BLOCKED computation_ended : queue;
    TERMINATED terminated;
  }
  ordering_criteria = {
```

}

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  process = {
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   time wcet;
    time current_deadline;
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  states = {
    RUNNING running : process;
    READY ready : select queue;
    BLOCKED blocked : queue;
    BLOCKED computation_ended : queue;
    TERMINATED terminated;
  }
  ordering_criteria = { lowest current_deadline }
```

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  process = {
    time period;
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  states = {
    RUNNING running : process;
    READY ready : select queue;
    BLOCKED blocked : queue;
    BLOCKED computation_ended : queue;
    TERMINATED terminated;
  }
  ordering_criteria = { lowest current_deadline }
  handler (event e) { ... }
}
```

Bossa event handlers

```
handler (event e) {
    On block.*
```

On unblock.preemptive

On bossa.schedule

Bossa event handlers

```
handler (event e) {
  On block.* { e.target => blocked; }
  On unblock.preemptive {
    e.target => ready;
    if (!empty(running) && e.target > running) {
      running => ready;
   }
  }
  On bossa.schedule {
    select() => running;
  }
  . . .
}
```

Other features

Timers

Used in EDF to wake a process at the beginning of its period.

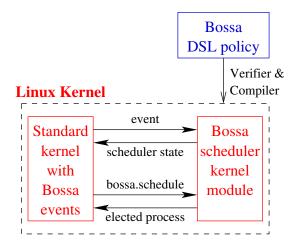
Interface functions

▶ Used in EDF to pause a process at the end of its computation.

Attach/detach functions

Used in EDF to check schedulability.

Preparing Linux for use with Bossa



Problem: adding Bossa event notifications to Linux

Traditional approach: Modify code and create a patch file.

- Tedious and error-prone.
- Only applies to one version of the OS.
- Non-modular.

Problem: adding Bossa event notifications to Linux

Our approach:

Aspect-Oriented Programming (AOP)

- Where: e.g., around call to try_to_wakeup.
- What: e.g., call Bossa event rts_unblock.
- Independent of line numbers and code details.
- Portable across multiple versions.

Components

- Describe the interface between the OS and the Bossa policy.
- Interface augmented with aspect rules.

Limitations of traditional aspect systems

Traditional aspect systems:

- > Put code before, after, and around functions and variables.
- ► No mechanism for referring to code sequences.

```
Linux schedule function:
```

```
schedule (void) {
  spin_lock_irq(&runqueue_lock);
  ... process election ...
  spin_unlock_irq(&runqueue_lock);
  ... context switch ...
}
```

Limitations of traditional aspect systems

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Linux schedule function:
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}
```

Extend AOP with Temporal Logic to describe code sequences.

Implemented using CIL and applied to Linux 2.4.18 and 2.4.24.

Verification

A Bossa policy must respect the scheduling requirements of the target OS.

Event types:

- Model of OS behavior.
- Created by the OS expert who integrates Bossa event notifications into the OS.
- > Example: unblock.preemptive for Linux.
 [tgt in BLOCKED] -> [tgt in READY]
 [p in RUNNING, tgt in BLOCKED] -> [[p,tgt] in READY]
 [tgt in RUNNING] -> []
 [tgt in READY] -> []

Event type checking

```
unblock.preemptive for Linux.
[tgt in BLOCKED] -> [tgt in READY]
[p in RUNNING, tgt in BLOCKED] -> [[p,tgt] in READY]
[tgt in RUNNING] -> []
[tgt in READY] -> []
```

Example definition:

```
On unblock.preemptive {
    e.target => ready;
    if (!empty(running) && e.target > running) {
        running => ready;
    }
}
Incorrect for linux!
```

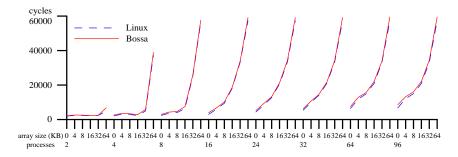
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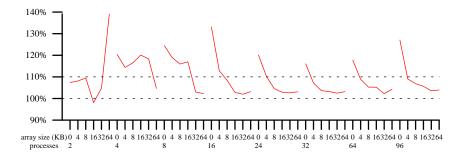
Example definition:

```
On unblock.preemptive {
    if (e.target in blocked) {
      e.target => ready;
      if (!empty(running) && e.target > running) {
        running => ready;
      }
    }
}
```

Performance: lat_ctx, context switch time



Performance: lat_ctx, increase as compared to Linux



Conclusions

Specialized process schedulers needed for demanding applications, but schedulers are not easy to implement in existing OSes.

Bossa provides:

- A DSL to ease the programming of scheduling policies.
- A Run-Time System implementing a scheduling interface.
- Verifications checking that a scheduling policy meets OS requirements.

Availability:

- Several versions of Linux 2.4.
- Teaching lab, based on Knoppix.

Ongoing work

- ▶ Modular Bossa [GPCE05], with a modular type system.
- Development of verified schedulers within the B framework.
- BossaBox: PVR with programmable scheduling.
- Coccinelle: automated support for device driver evolution [ACP4IS].

More information

- Framework [HASE 2005]
- Aspects and components [SIGOPS 2004]
- Verification [GPCE 2004]
- Generalization to a scheduler hierarchy [PEPM 2004]

http://www.emn.fr/x-info/bossa/