## **Multihoming: An Overview** & a brief introduction to GSE(8+8)

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- Get a slice of address from your ISP
- There is only one way to send, and one way to receive data
- ISPs inject aggregated prefixes into global routing
- Life is simple

## **The Internet is made of multiple ISPs !**



- Connect to multiple ISPs!
  - Get the best from each and all
- Bringing more questions to life ...



Multihoming

## **Multihoming: Sending Data**



- Which way to get out?
  - You probably pay different price for different ISPs
  - The best choice may depend on specific destinations you are sending to
  - If more than one exit router: decision is made *inside the site*
- Outbound traffic engineering

## **Multihoming: Receiving Data**



- You'd like
  - Traffic coming from a cheap link as long as it works, or
  - utilizing all links in certain proportion
- Inbound traffic engineering

### **Multihoming: lets not forget the middle**



- When a *destination* is more than one AS hop away, a transit ISP may wish to know whether there are multiple ways to reach , so that it can choose the *"best"* one
- *Transit traffic engineering* <sup>™</sup>

## **Multihoming: Addressing**



For a multihomed site:

- Where does one get the address?
- What/which address to use for
  - source address?
  - destination address?

#### **Multihoming Address: PI Prefix**



- (some)sites can get a prefix allocation from RIRs directly
  - *Every* PI prefix adds an entry into the global routing table

Global Routing Table

0.1.10.0/20

#### **Multihoming Address: PI with TE**



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#### **Multihoming Address: PA**



2.2.0.0/16

2.2.0.0/16

2.2.8.0/20

#### **Multihoming Address: PA + hole + split**



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# **Multihoming: Today's Practice**

- Use of PI prefixes, or punching holes on PA prefixes
- Works as far as TE is concerned (only)

these days, many people take TE to mean "break up my portable address block in small parts". That is one way to do it, and an effective one, but also the least scalable one -- Iljitsch van Beijnum 2/24/06

- Lots worries about impact on routing scalability
  - Goes up with the number of multihomed sites
  - Goes up further with additional prefix split for inbound TE

## **One of Earlier Proposals: GSE**

• GSE: Global, Site, and End-system address elements, proposed by Mike O'Dell in 1996-97 <u>http://www.watersprings.org/pub/id/draft-ietf-ipngwg-gseaddr-00.txt</u>

Transit backbone

• The basic idea:

- Separate public and private topologies
- *Insulate* customer site from the global provider topology

customer sites

## **How GSE Works**

Proposed IPv6 address structure:



- Internal packet delivery without RGs
- External packet delivery: defer/hide RG AMAP
- Multihomed sites get multiple RGs

## How GSE Works: more detail

- For outbound traffic:
  - Get destination address and RG from DNS lookup
  - Put on source RG when packets exiting local site



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- For outbound traffic:
  - Get destination address and RG from DNS lookup
  - Put on source RG when packets *exiting* local site
- For inbound traffic
  - Take off destination RG at entrance to destination site
  - Keep source RG for returning traffic



## What Problems GSE Solved

- Making customer sites unaware of the transit backbone or provider change
   Eliminate renumbering caused by change of providers
- Providing freedom for performing aggregation as needed in the provider space

## What Issues Left Open

- *1. Inbound traffic engineering*: which destination RG to put on each packet?
- 2. *Outbound traffic engineering*: How to select exit router?
- *3. Transit traffic engineering*: How to tell whether packets carrying different destination RGs belong/not belong to the same destination site?



## What New Issues Introduced

- Interactions with DNS: since one learns destination RG from DNS lookup
  - Would the RGs for DNS root servers (ever) change?
  - RGs for other DNS servers *will* change
    - thorough analysis needed to understand the implication and impact
- What if remote link of the selected destination RG fails?
  It is proposed *not* to use dynamic DNS for link status changes
- For IP Tunnels across RG boundaries:
  - What source/destination RGs going into the tunnel?
  - How to handle the RGs when packets get out (land in a different site)?
- Would allowing address rewriting make TCP connection hijacking much easier than today?

## **Another Early Proposal: Map & Encap**

- RFC1955 (June 1996) by Bob Hinden
- Basic idea: Putting ISPs in a separate address space from customers
- Benefit & issues: similar to GSE



# Yet Another Early Proposal: Metro-based address

- (expired) Internet-Draft by Deering & Hinden: "IPv6 Metro Addressing" March 1996
- Basic idea: address allocation by metro areas
- Benefit: provider-independent addressing
  - Support for multihoming
  - Elimination of renumbering when changing providers
- Issues: provider-independent addressing
  - Requiring providers to inter-connect in all metro areas
  - Lack routing policy support

# **Multihoming: Summary of Issues**

In today's practice:

- PI prefix is effective for multihoming and TE, but raises scaling concerns
- Use of multiple PA prefixes bring up new issues
  - Source address selection
  - Source exit router selection (outbound traffic engineering)
  - Destination address selection (inbound traffic engineering)
  - Who/where/how to control the above
  - Transit traffic engineering
- A few related issues not discussed
- E.g. ingress filtering

# **Departing Words**

- GSE brought up a new approach to the problem
  - An interesting, but early stage, proposal; a number of important issues remain open
- (Quote from NANOG35 IAB BOF) "routing has always been working with prefixes, which are *locators*"
  - locator/identifier overload/split is an issue for transport and above, but *not* the cause of today's routing scalability problem
- (Tony Li) creating a scalable routing subsystem is paramount, as without that, we effectively have no network.
- The community must work together, *step by step*, towards a solution



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