

Implementing Direct Anonymous Attestation for the TPM Emulator Project

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- 1 Introduction
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“THE **TRUSTED COMPUTING GROUP (TCG)** IS AN INDUSTRY ORGANIZATION THAT AIMS TO DEVELOP AND PROMOTE OPEN, VENDOR-NEUTRAL STANDARD SPECIFICATIONS FOR TRUSTED COMPUTING BUILDING BLOCKS AND SOFTWARE INTERFACES ACROSS MULTIPLE PLATFORMS.” [TCG]

- Hardware: Specification of the **Trusted Platform Module (TPM)**
 - Version 1.1b (February 2002), Version 1.2 (Revision 62: October 2003, Revision 85: February 2005, Revision 94: March 2006)
- Mainboard/Firmware/BIOS: **Platform Specific Specifications**
 - PC Client TPM Interface Specification (TIS), Version 1.2
 - Implementation Specification for Conventional BIOS, Version 1.2
- Software: Specification of the **TCG Software Stack (TSS)**
 - `TSS ::= (TCG | TPM) Software (Stack | Specification)`
 - Version 1.1 (September 2003), Version 1.2 (January 2006)
 - News: Auditing, delegation, monotonic counters, DAA, ...
- Network: Specification of the **Trusted Network Connect (TNC)**

Main Functionality of the TPM:

- “Cryptographic co-processor” (RNG, SHA-1, HMAC, RSA)
- Hardware protected storage for cryptographic keys
- Measurement of the platform configuration (PCR)
 - Application: Secure Boot and **Remote Attestation**
- Sealed Storage: Decryption keys are tied to a PCR value

Remote Attestation: Certify a platform characteristic (e.g. the current software configuration by means of PCR values) to a remote party.

Application:

- System integrity check (to detect corrupted software, e.g. root kits)
- Enforcement of a corporate-wide software stack
- Digital rights management (DRM), pay-per-use services
- Product activation, tethering, and customization
- Vendor lock-in, forced upgrades and downgrades, ...

Previous Solution (Trusted Third Party):

introduced by TPM Specification, Version 1.1b

- 1 Owner initiates the creation of an attestation identity key (AIK)
 - AIKs are special-purpose signature keys (non-migratable, RSA 2048 bit, $e = 2^{16} + 1$) generated and protected by a TPM
- 2 Privacy-CA certifies this AIK, if the platform is able to show its conformance with a policy (e.g. valid EK/platform credentials)
- 3 AIK and the obtained AIK-certificate are used by the platform to perform a desired remote attestation (i.e. sign PCR values)

Disadvantages:

- Different attestations are **linkable**, if the same AIK is used multiple times. Thus owners should always create fresh AIKs.
- The Privacy-CA is a very sensitive entity. Therefore it must be **carefully protected** and maintained to guarantee security.
- The Privacy-CA must be **highly available**, because it is involved in every attestation. (but this contradicts the point above)

Brickell, Camenisch, and Chen [BCC04]: **Direct Anonymous Attestation**

- Combine ideas from group signature schemes with the efficient Camenisch-Lysyanskaya anonymous credential system [CL01, CL02]
- “GROUP SIGNATURE SCHEME WITHOUT ANONYMITY REVOCATION”
- DAA entities:
 - Host/TPM:** TC-platform consists of a host and a trusted observer (TPM)
 - Issuer** issues a DAA-certificate, if TC-platform possess a valid EK (policy)
 - Verifier** provides a service, if DAA-signature is valid w.r.t. desired message
- DAA sub-protocols:
 - DAA-Join:** TPM chooses a secret f and performs a two-party protocol with an issuer to obtain a secret DAA-certificate $A_I(f)$ (i.e. a CL-signature on f).
 - DAA-Sign:** TPM signs an attestation message m (e.g. hash value of an AIK). The appended non-interactive zero-knowledge proof of knowledge shows that the corresponding DAA-signature $\sigma_I(f, m)$ is valid w.r.t. f , m , and I (issuer key).
- Issuer and verifier are able to detect “broken TPMs” (Rogue Tagging)
- Unforgeability of DAA-certificates relies on **Strong RSA Assumption**
- Unlinkability of DAA-certificates/signatures relies on **DDH Assumption**

Advantages:

- DAA-certificates need to be **issued only once** (no bottleneck)
- Issuer and verifier **cannot link** DAA-certificates and DAA-signatures, even if they are the same entity (“repairs the broken business model”)
- Anonymity **degradation** is possible (named base vs. random base)
 - Detect and exclude malicious TPMs (e.g. black list)
 - Perform frequency analysis (DoS attack on issuer/verifier)

1st Extension: (does not break the current TPM Specification)

- Better privacy by combining DAA with a Privacy-CA which issues “one-time” DAA-certificates (Camenisch, ESORICS 2004)

2nd Extension: (breaks the current TPM Specification)

- Ensure anonymity on **untrusted** TPMs (Camenisch, unpublished)

DAA in TCG Specifications: introduced by TPM and TSS Specification, Version 1.2

TPM 1.2 TPM part of DAA-Join and DAA-Sign (structures and commands)

TSS 1.2 Host part of DAA-Join and DAA-Sign, NIZK verification (issuer key), supplemental functions for issuer and verifier, . . .

“Additional Features” Introduced by TCG Specifications:

TSS 1.2 **Arbitrary attributes** (e.g. expiration date) for DAA-certificates

TSS 1.2 Optional **anonymity revocation** based on verifiable encryption (cf. Camenisch and Shoup, CRYPTO 2003)

Technical Problems:

- DAA-Join resp. DAA-Sign are highly resource intensive protocols (TPM: 11 resp. 7 modular exponentiations with large exponents)
- `TPM_DAA_Join` and `TPM_DAA_Sign` are executed in atomic stages; in-between they may be interruptible by other commands (save context)
- DAA-Join must not run arbitrarily interleaved (restriction handled by TSS)
- Protection against timing attacks? (e.g. exp. in stage 4 of `TPM_DAA_Join`)

Mario Strasser: Software-based TPM Emulator for Linux [St04]

<https://developer.berlios.de/projects/tpm-emulator/>

Goal: Create a fully working Trusted Platform Module emulator according to TCG Specification, Version 1.2 (Revision 62 → 94)

Application: Explore TPMs for educational/experimental purposes

Current State: Release 0.3 (January, 2006), GNU GPL v2

- Kernel module `tpm_emulator.ko` (provides char. device `/dev/tpm`)
- Currently, 80 out of 120 TPM commands are implemented (admin startup, admin testing, admin opt-in, admin ownership, auditing, storage functions, cryptographic functions, endorsement key handling, identity creation, integrity collection and reporting, authorization sessions, session management, eviction, timing ticks, transport sessions, monotonic counter, DAA, deprecated commands)
- Not yet/only partially implemented: capability, migration, maintenance, identity activation, delegation, NV storage
- Packetized for Gentoo Linux (`$ emerge tpm-emulator`)

Prerequisites:

- Linux Kernel 2.6.x, GNU Compiler Collection, ...
- GNU Multiple Precision Arithmetic Library (libGMP)

Roadmap/TODO:

- 1 Conformance with Revision 94 of the TPM Specification 1.2
- 2 Obtain better portability (kernel space vs. user space)
1st Problem: Kernel stack size is very limited
(architecture dependent, e.g. 4K resp. 8K on x86)
2nd Problem: Persistent storage is needed to save the state

Possible Solution:

- Dummy “hardware interface” in the common TPM device driver
 - TPM emulator serves only as user space daemon
- 3 Implementation of all mandatory commands (v1.2 rev 94)
 - 4 Adding optional commands and algorithms (e.g. AES)

- Approximately 3850 lines of code (including 600 lines of comments)
- Implementation time: \approx 9 days (3–4 hours per day)
- Testing time: \approx 6 weeks (IBM DAA Test Suite [Zi05])

Requirements:

- Kernel stack size of at least 8K (libGMP calls, large structures)

Implementation Problems:

- Missing kernel debugger (e.g. to detect call paths with stack overflows)
- Alignment of large integers in combination with hashing
- TPM specification contains many typographical errors:
 - “TPM computes a TPM-specific secret f_0 (104-bit) = $f \bmod 2^{104}$ ” (Part 1, rev 94)
 - “#define DAA_SIZE_r3 158” (Part 2, rev 62) “#define DAA_SIZE_r3 168” (Part 2, rev 85)
 - “obtain DAA_SIZE_NT bits from RNG” (20 bits vs. 20 bytes) (Part 3, rev 94)

- The implementation works well and was carefully tested with the IBM DAA Test Suite [Zi05] (thanks to Roger Zimmermann for his support).
- TPM and TSS 1.2 specifications (even the current revision 94) contain many typographical errors, often change data structures, and thus are **difficult to implement** resp. **hard to keep up to date**.
- Open-source TSS: Kent Yoder (IBM, TrouSerS Project [Yo06])
 - “RIGHT NOW WE HAVE TWO PEOPLE IMPLEMENTING DAA.”
 - “WE’RE PLANNING ON INTEGRATING THE IMPLEMENTATION SOME TIME IN THE NEXT COUPLE MONTHS.”
- **Contributions to TPM-Emulator Project [St06] are very welcome!**

Thank you!

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