

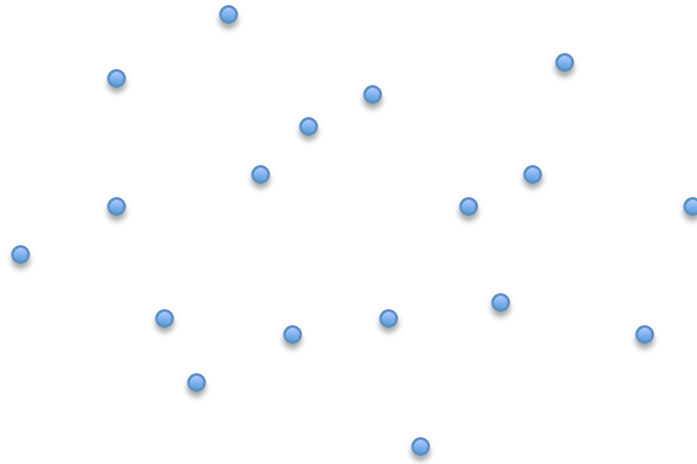
# Convex Hull

(divide and conquer)

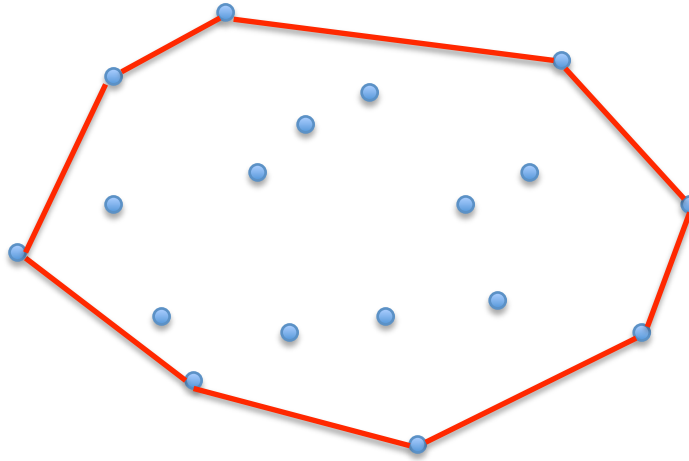
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Grenoble-INP / april 2010

# Set of points $Q$



# Convex hull (definition)

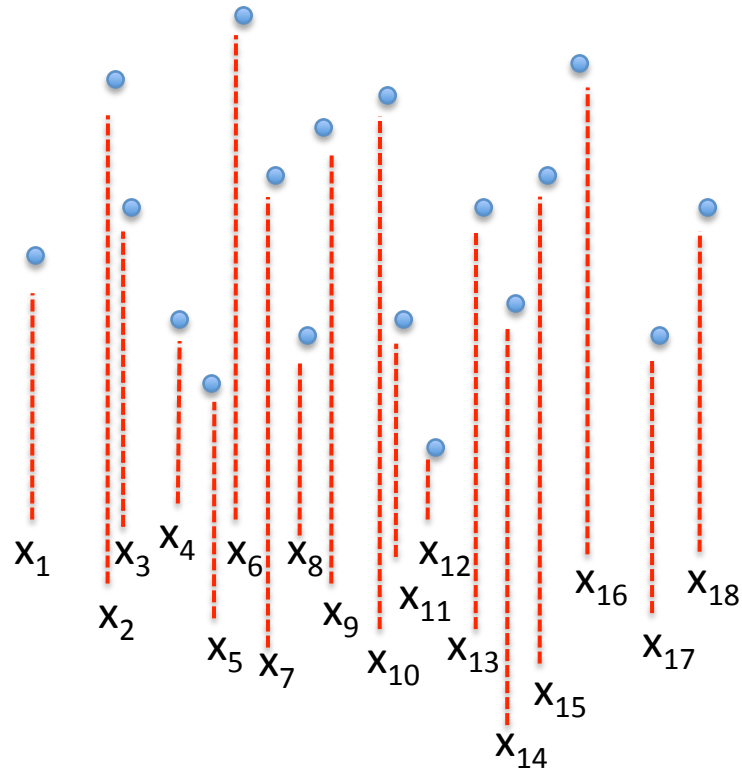


H is the smallest convex polygon that contains all the points of Q

# Principle

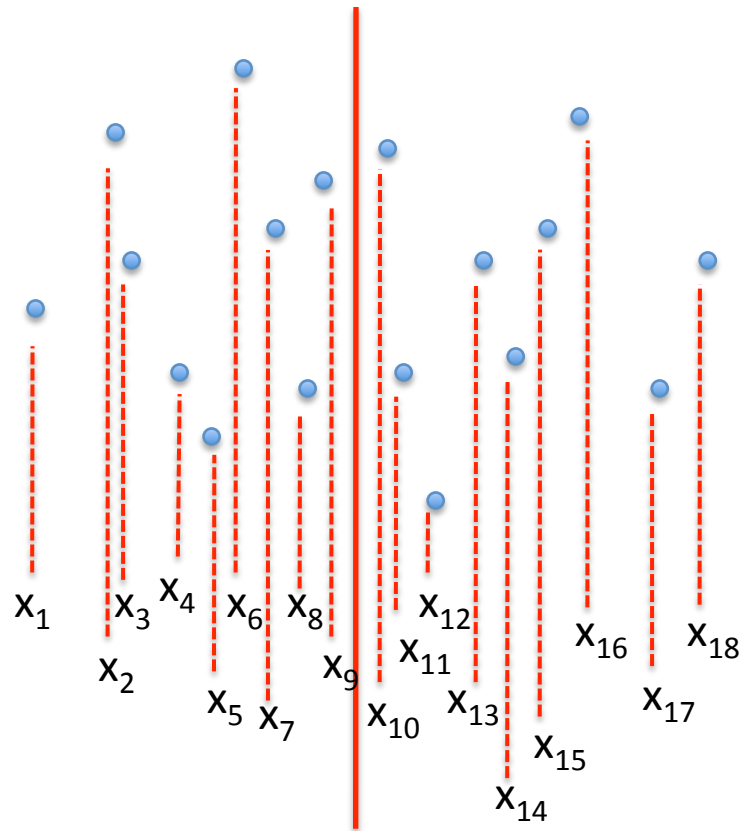
Decompose the set of points in equal parts ( $Q_{\text{left}}$  and  $Q_{\text{right}}$ )  
Solve the sub-problems respectively on  $Q_{\text{left}}$  and  $Q_{\text{right}}$   
Merge both convex hulls  $H_{\text{left}}$  and  $H_{\text{right}}$

# 1. decomposition



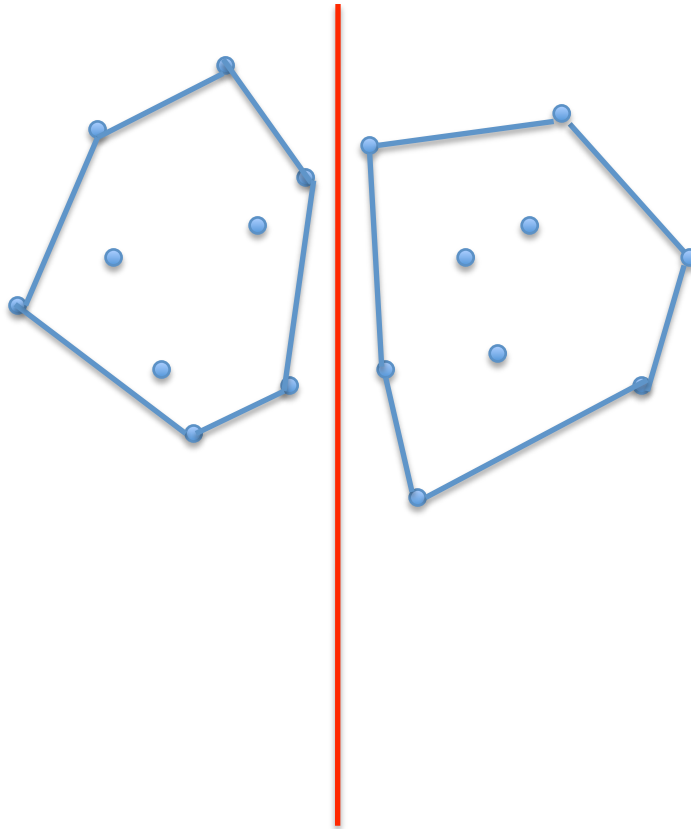
Sort the points by increasing abscissa

# 1. decomposition



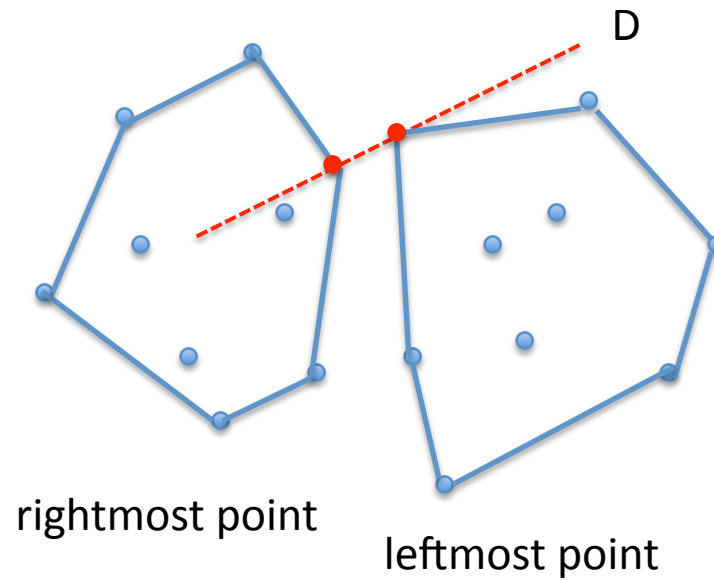
Split the points into two sets of equal size

## 2. Solve sub-problems



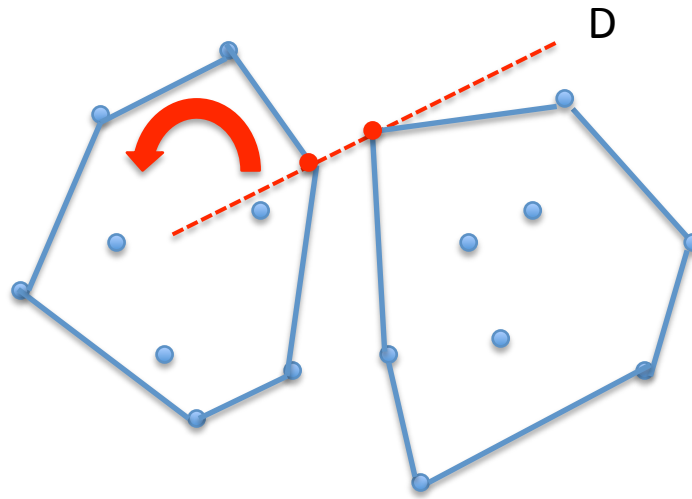
Compute the convex hulls on Qleft and Qright

### 3. Merge both solutions



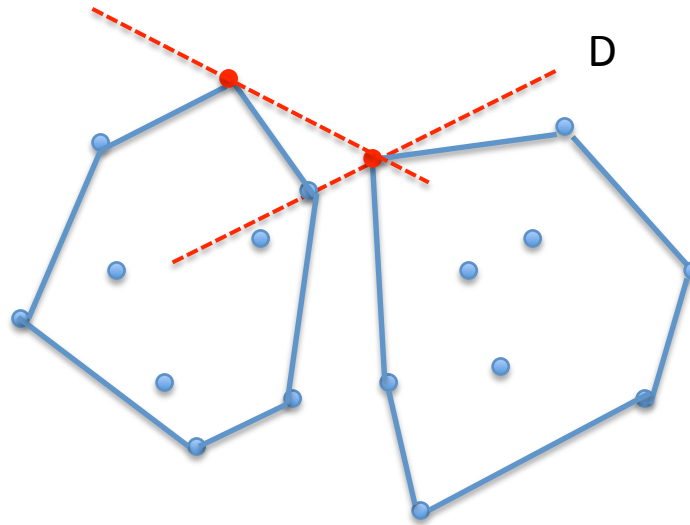


### 3. Merge both solutions



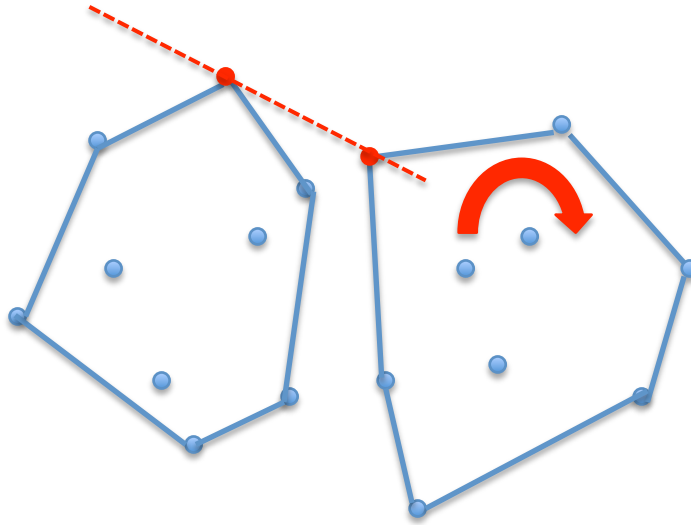
Determine the upper tangente

### 3. Merge both solutions



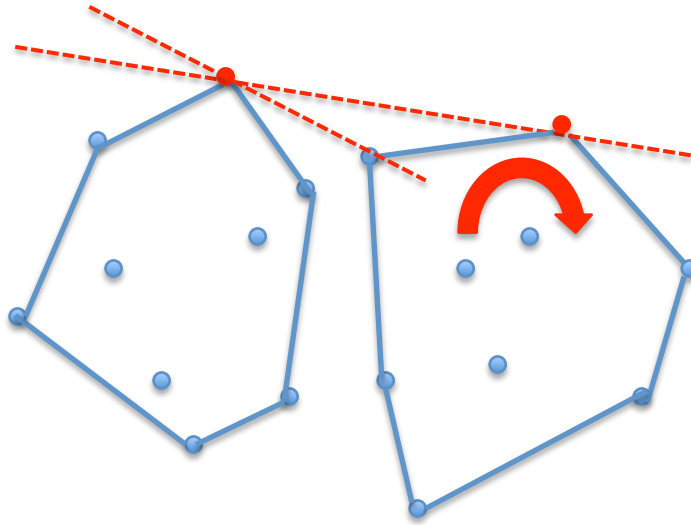
Determine the upper tangente

### 3. Merge both solutions



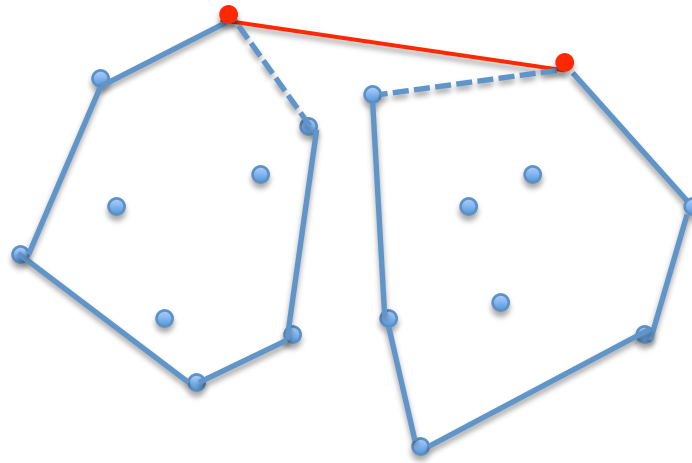
Determine the upper tangente

### 3. Merge both solutions



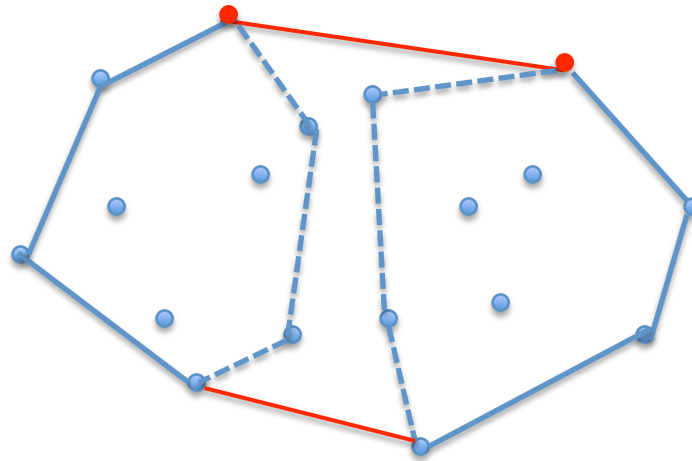
Determine the upper tangente

### 3. Merge both solutions



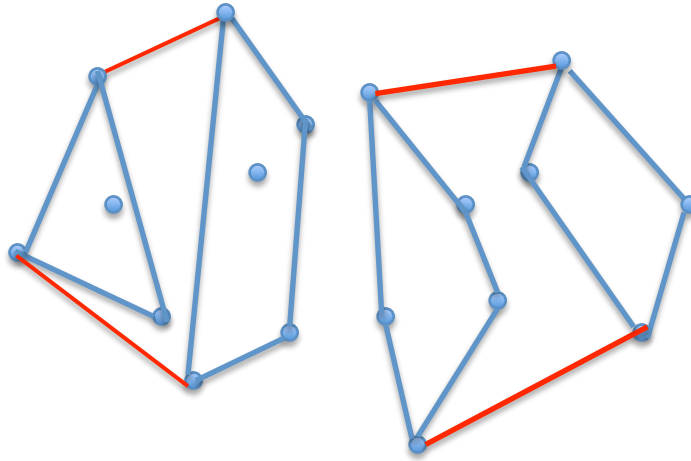
Determine the upper tangente

### 3. Merge both solutions

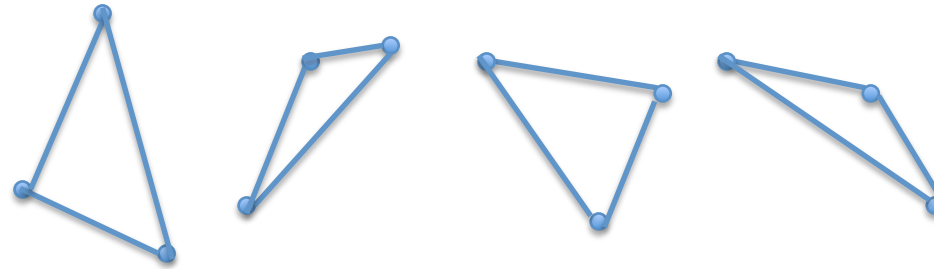


Determine the lower tangente

# Recursive decomposition



# Basis of the induction by « brute force » case analysis



4 types of triangles



Cost analysis:  $T(n) = 2 T(n/2) + O(n) = O(n \cdot \log_2(n))$