Tuning & Applying Best Practices On Very Large Databases (VLDB)

Tips, tricks & success stories from real life

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Agenda

- Real-life scenario 1 How a daily batch job performance consistency was achieved
- Real-life scenario 2 ETL data loading and BI queries dramatic improvements
- Best Practices for configuration, sizing, performance, storage consideration
 - Golden Rules for VLDB configuration
 - Backup & Recovery best practices
 - Optimizer Statistics management
 - ASM best practices
 - Performance optimization
 - Exadata solutions
 - Active Data Guard solutions

Not a self marketing...

A tiny inspiration

Over 20 years of IT experience 16+ years as an Oracle DBA Oracle ACE Director Oracle 10g Certified Master(OCM) Oracle 10g RAC Certified Expert OCP v8i,9i,10g & 11g ITIL v3 Foundation Certified Oracle Database 12c <u>beta tester</u>

SNC ID: @sjaffarhussain http://jaffardba.blogspot.com Certified Professional

ORACLE 10g Certified Master

Certified Expert

10g RAC Administrator





TID-BOLODISTOFTIC FEAR

Syed Jaffer Hussain 🛤

The learning name stops for Grocke RAC expert.



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Oracle Scientific collaboration Technologies

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Products, "and later indexity, and and going new surface ingeneration and going respect on the copyrights."



Technologist of the Year, DBA 2011

http://www.oracle.com/technetwork/issue-archive/2012/12-jan/o12awards-tech-1403083.html

Co-authored ...



Oracle 11g R1/R2 Real Application Clusters Essentials

Design, implement, and support complex Oracle 11g RAC environments for real-world deployments

Ben Prusinski Syed Jaffer Hussain Converced Natures (PACKT) enterprise ⁸⁸

Expert Oracle RAC 12c

GAIN DEEP EXPERTISE IN MANAGING ORACLE REAL APPLICATION CLUSTERS

Apress[®]

Syed Jaffar Hussain, Tariq Farooq, Riyaj Shamsudeen, and Kai Yu

Cupyrighted Material

Presented by : Syed Jaffer Hussain

Up coming ...

Comprehensive, end-to-end insight into Oracle's flagship database machine



Oracle Exadata

Tariq Farooq | Charles Kim | Nitin Vengurlekar Sridhar Avantsa | Guy Harrison | Syed Jaffar Hussain



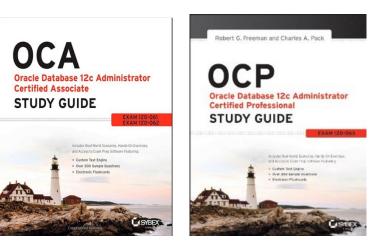




Oracle Data Guard 11gR2 Learn how to build and maintain Data Guard configurations

Beginner's Guide

Technical Reviewer ...



VLDB - Definition

Data Warehouse -> VLDB -> Big Data

- Holds extremely high quantity of data
- Occupies large volume of storage
- Contains historical data for decision making
- Often updated through ETL tools
- A.k.a Big Data
- DW as a foundation for Big Data

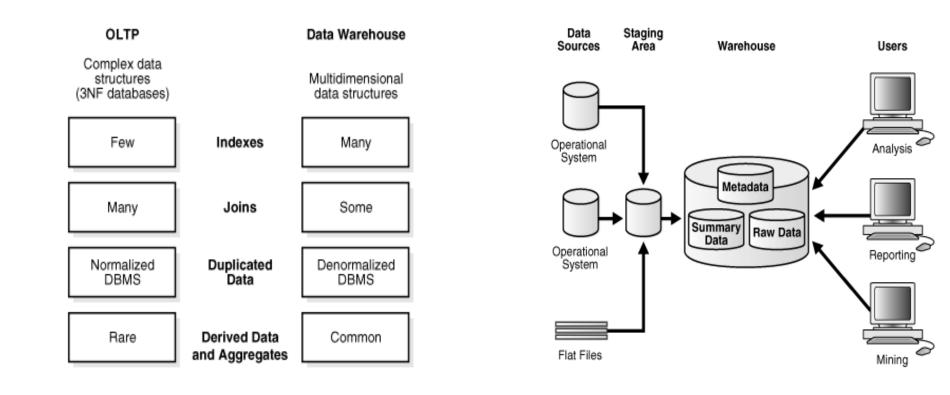
VLDB – Key challenges

- Configuration
- Storage
- Performance
- Maintenance
- Administration
- Availability
- Server resources

Data Warehousing Concepts - Overview

Contrasting OLTP and Data Warehousing Environments

Data Warehouse Architecture: with a Staging Area



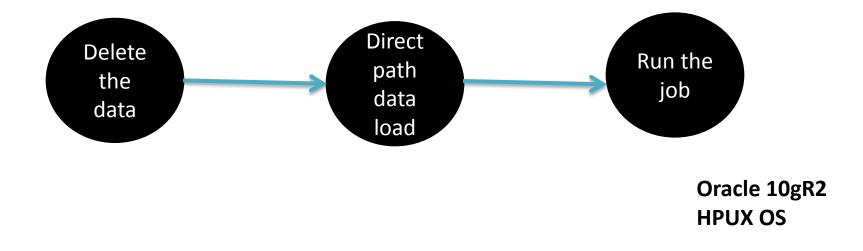
Scenario 1

How a daily batch job performance consistency was achieved

Scenario 1

How a daily batch job performance consistency was achieved

Reconciliation batch job cycle



Scenario 1

How a daily batch job performance consistency was achieved

Major obstacles

- Job completion inconsistency
- Taking 2 hours to 24 hrs to complete
- Impact on the job dependency



Scenario 1

How a daily batch job performance consistency was achieved

Observation

- Consistent amount of data growth each day
- Queries favoring NESTED LOOP



Scenario 1

How a daily batch job performance consistency was achieved

Temporary fix

- Table and schema re-org
- Collecting stats
- Using Hints (HASH JOIN)
- Migration to other database



Scenario 1

How a daily batch job performance consistency was achieved

Remedy

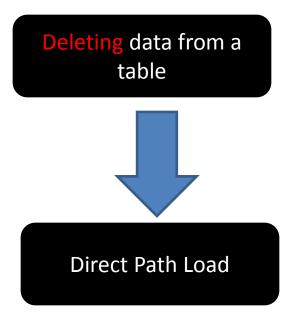
A <u>little ADJUSTMENT to the CODE</u> did the MIRACLE!



Scenario 1

How a daily batch job performance consistency was achieved

Code review and Findings





Scenario 1

How a daily batch job performance consistency was achieved

> Delete was REPLACED with truncate



Scenario 1

How a daily batch job performance consistency was achieved

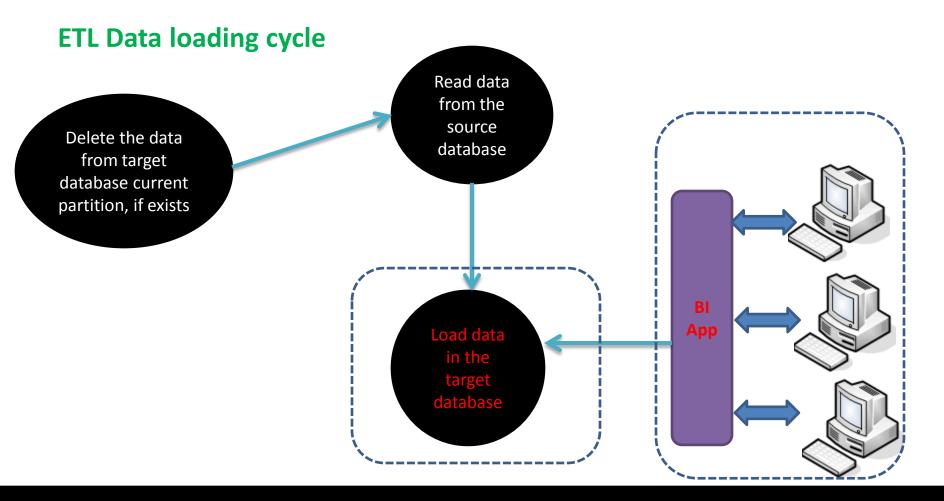
- Its been over 6 yrs or so, the JOB is not taking more than 45 min. to finish!
- > No more fragmentation
- > Queries opting HASH JOIN and finishing on TIME

Scenario 2

ELT data loading and BI queries dramatic improvements

Scenario 2

ELT data loading and BI queries dramatic improvements



Scenario 2

ELT data loading and BI queries dramatic improvements

ETL Data loading cycle - statistics

- ✓ Over 1 millions are inserted a table using the plain inserts
- ✓ Monthly partitioned table
- ✓ Over 8 indexes
- \checkmark 6 tables involved in the process

Scenario 2

ELT data loading and BI queries dramatic improvements

ETL Data loading cycle - Concerns

- ✓ Through put, <150 records per second
- ✓ Taking 4 hours to complete the load job
- ✓ Stats gather was taking 2-4 hours
- ✓ BI queries on the tables were taking over 2 hours
- ✓ Affecting subsequent job cycle, impacting other databases and busiess
- ✓ Business users complains, escalations if the reports did not deliver by 7.30am

Scenario 2

ELT data loading and BI queries dramatic improvements

Action force committee

- ✓ Emergency meeting was called
- Discussed the seriousness of the issues
- ✓ Come-up with the list of action plans (around 27 actions)
- ✓ Targets/deadlines were set



Scenario 2

ELT data loading and BI queries dramatic improvements

Database action plan

- 1. Monitor index usage to drop unwanted indexes
- 2. Verify chained rows
- 3. Splitting monthly partitions to weekly
- 4. Incremental gather stats
- 5. Increase redo log size
- 6. Enable Parallelism
- 7. Convert Global indexes to Local indexes
- 8. Adding resources, CPU & Memory to the database server
- 9. Increase SGA
- 10. Rewriting some of the queries(logic)
- 11. Enable Temporary Tablespace Group Management
- 12. Offloading queries to standby database
- 13. Restructured the partition key on the tables

Scenario 2

ELT data loading and BI queries dramatic improvements

Get rid-off unwanted indexes

- ✓ Index usage monitoring / dropping unwanted indexes
 - ✓ Managed to reduce number of indexes, dropped and composite indexes
 - ✓ Made the indexes invisible before drop

NDEX_NAME	TABLE_NAME	MON USE START_MONITORING	END_MONITORING
DRILL T24_CONTRACT_DOA_IDX	DRILL_T24_CONTRACT_BALANCE	YES NO 12/11/2013 12:08:5	8
DRILL_T24_CONTRACT_N1	DRILL T24 CONTRACT	YES YES 12/11/2013 10:17:2	3
DRILL_T24_CONTRACT_N2	DRILL_T24_CONTRACT	YES NO 12/11/2013 10:17:2	3
DRILL_T24_CONTRACT_N3	DRILL_T24_CONTRACT	YES NO 12/11/2013 10:17:2	3
DRILL_T24_CONTRACT_N4	DRILL_T24_CONTRACT	YES YES 12/11/2013 10:17:2	3
DRILL_T24_CONTRACT_OPENING_N1			
DRILL_T24_CONTRACT_OPENING_N2		YES NO 12/11/2013 12:08:5	
DRILL_T24_CONTRACT_OPENING_N3	DRILL_T24_CONTRACT_BALANCE	YES YES 12/11/2013 12:08:5	8
DRILL_T24_CONTRACT_OPENING_N4			
DRILL_T24_CONTRACT_OPENING_N5	DRILL_T24_CONTRACT_BALANCE	YES YES 12/11/2013 12:08:5	9
DRILL_T24_CONTRACT_OPENING_PK	DRILL_T24_CONTRACT_BALANCE	YES NO 12/11/2013 12:08:5	9
DRILL_T24_CONTRACT_PK DRILL_T24_TRANSACTION_N1 DRILL_T24_TRANSACTION_N5	DRILL_T24_CONTRACT	YES YES 12/11/2013 10:17:2	3
DRILL_T24_TRANSACTION_N1	DRILL_T24_TRANSACTION	YES YES 12/11/2013 10:17:5	6
DRILL_T24_TRANSACTION_N5	DRILL_T24_TRANSACTION	YES YES 12/11/2013 10:17:5	6
DRILL_T24_TRANSACTION_N6	DRILL_T24_TRANSACTION	YES YES 12/11/2013 10:17:5	6
DRILL T24 TRANSACTION N7	DRILL T24 TRANSACTION	YES YES 12/11/2013 10:17:5	6
DRILL_T24_TRANSACTION_PK	DRILL_T24_TRANSACTION	YES NO 12/11/2013 10:17:5	6

Scenario 2

ELT data loading and BI queries dramatic improvements

Chained rows

- ✓ Verify chained rows
 - ✓ Over 60 thousand chained rows
 - ✓ The table had 1% of row chaining
 - ✓ Migrated to a higher block tablespace

Scenario 2

ELT data loading and BI queries dramatic improvements

Database action plan

- ✓ Splitting monthly partitions to weekly with auto partitioning option
- ✓ Change in partition column
- ✓ Incremental gather stats
- ✓ Increase redo log size from 100MB to 500MB
- ✓ Convert Global indexes to Local indexes
- ✓ Adding resources, CPU & Memory to the database server
- ✓ Increase SGA
 - ✓ Incremental stats : from 2-4 hours to <30min
 - ✓ Inserts : increased from <150 RPS To >2000 RPS
 - ✓ The insert job completes in <6 min in contrast to 4 hrs earlier
 - ✓ BI reports improved from 18 min to 27 seconds

Scenario 2

ELT data loading and BI queries dramatic improvements

Database action plan

- ✓ Rewriting some of the queries(logic)
 - ✓ 5hrs job was brought down to 3 min

Scenario 2

ELT data loading and BI queries dramatic improvements

Testimonials

The Query which used to take 18 minutes is now taking only 27 seconds to get the data from the Database. Below are the snapshots: 🛃 Track 🔹 🐺 Drill 🔹 🌾 Filter Bar 🔛 Freeze 🔹 Web Intelligence 🔹 📋 🧀 🎽 📲 🔹 🏭 🔹 9 CH 20 -Del Outline Navigation Map 1 Line of Business (All values) 💉 Segment Description (All values) 💙 Segment Statistics (Partner Numbers & Deposits, by Segme All Retail Partners All Retail Partners Across Regions Segment Statistics All Retail Partners Across Branches NO Partner Numbers & Deposits, by Segment Refresh Date: 21-JAN-2014 04:59:17 (?)As of Date: Numb Line of Business Segment Description Acco Large Corporate Medium Corporate **Micro Enterprise** Corporate Sr **Refreshing Data** - N Line of Business Af Last refresh time: Oh Om 27s

Scenario 2

ELT data loading and BI queries dramatic improvements

Database action plan

\checkmark Restructured the partition key on the tables

SELECT FROM_DATE , TO_DATE, ACCOUNT_NUMBER FROM TABLE_BEFORE_MODIFIATION WHERE '01-DEC-2012' BETWEEN FROM_DATE AND TO_DATE

SELECT STATEMENT,			Cost= <mark>17373</mark>	Cardinality=2089220	Bytes=64765820	Time=244	IO cost=17264	CPU cost=1297607466
GOAL = ALL ROWS								
PARTITION RANGE			Cost=17373	Cardinality=2089220	Bytes=64765820	Time=244	IO cost=17264	CPU cost=1297607466
ITERATOR								
TABLE ACCESS FULL	Object	Object name=	Cost=17373	Cardinality=2089220	Bytes=64765820	Time=244	IO cost=17264	CPU cost=1297607466
	owner=xxxx	TABLE_BEFORE_MODIFICATIO						
		Ν						

SELECT FROM_DATE , TO_DATE, ACCOUNT_NUMBER

FROM TABLE_AFTER_MODIFICATION

WHERE '01-DEC-2012' between from_date AND TO_DATE

SELECT STATEMENT, GOAL = ALL ROWS			Cost= <mark>60</mark>	Cardinality=578568	Bytes=20249880	Time=1	IO cost=46	CPU cost=168083738
PARTITION RANGE ITERATOR			Cost=60	Cardinality=578568	Bytes=20249880	Time=1	IO cost=46	CPU cost=168083738
PARTITION HASH ALL			Cost=60	Cardinality=578568	Bytes=20249880	Time=1	IO cost=46	CPU cost=168083738
TABLE ACCESS FULL	Object owner= XXXX	Object name= TABLE_AFTER_MODIFICATION	Cost=60	Cardinality=578568	Bytes=20249880	Time=1	IO cost=46	CPU cost=168083738

		TABLE_BEFORE_MODIFICATION	TABLE AFTER MODIFICATION
PARTITION	RANGE	FROM_DATE	TO_DATE
SUB PARTITION	HASH(4)	N/A	FROM_DATE
INDEXES		BTREE GLOBAL INDEX	BTREE LOCAL INDEX

SELECT STATEMENT GOAL = ALL ROWS			Cost= <mark>3126</mark>	Cardinality=2072103	Bytes=64235193	Time=44	IO cost=3070	CPU cost=671176973
INDEX FAST FULL	Object	Object	Cost=3126	Cardinality=2072103	Bytes=64235193	Time=44	IO cost=3070	CPU cost=671176973
SCAN	owner=XXXX	name=TABLE_BEFORE_MODIFICATIO						
		Ν						

BTREE LOCAL INDEX ON FROM_DATE, TO_DATE ON AFTER_MODIFIATON_TABLE

SELECT STATEMENT, GOAL = ALL ROWS			Cost= <mark>103</mark>	Cardinality=821359	Bytes=28747565	Time=15	IO cost=1031	CPU cost=73807108
PARTITION RANGE			Cost=1037	Cardinality=821359	Bytes=28747565	Time=15	IO cost=1031	CPU cost=73807108
ITERATOR PARTITION HASH			Cost=1037	Cardinality=821359	Bytes=28747565	Time=15	IO cost=1031	CPU cost=73807108
ALL TABLE ACCESS BY	Object	Object name=	Cost=1037	Cardinality=821359	Bytes=28747565	Time=15	IO cost=1031	CPU cost=73807108
LOCAL INDEX	owner=XXXX	TABLE_AFTER_MODIFICATION	0050 1057		57105 207 17505	Time 15	10 0050 1051	
ROWID	Ohiost	Object come	Cast 000	Candinality, 7202		Time 14	10 east 002	CDU east 72201050
INDEX RANGE SCAN	Object owner=XXXX	Object name= INDEX_ON_AFTER_TABLE_MODI	Cost=988	Cardinality=7392		Time=14	IO cost=982	CPU cost=73391858

BTREE LOCAL INDEX ON TO_DATE, FROM_DATE

SELECT STATEMENT, GOAL = ALL ROWS			Cost= <mark>555</mark>	Cardinality=821359	Bytes=28747565	Time=8	IO cost=549	CPU cost=75584538
PARTITION RANGE ITERATOR			Cost=555	Cardinality=821359	Bytes=28747565	Time=8	IO cost=549	CPU cost=75584538
PARTITION HASH			Cost=555	Cardinality=821359	Bytes=28747565	Time=8	IO cost=549	CPU cost=75584538
TABLE ACCESS BY LOCAL INDEX ROWID	Object owner=OFDM	Object name= TABLE_AFTER_MODIFIATION	Cost=555	Cardinality=821359	Bytes=28747565	Time=8	IO cost=549	CPU cost=75584538
INDEX RANGE SCAN	Object owner=OFDM	Object name= INDEX_ON_AFTER_TAB_MODIFICATIO N	Cost=502	Cardinality=8142		Time=8	IO cost=496	CPU cost=75131718

Dears,

Please accept my sincere gratitude. The business are happy about the change in performance.

Well done Jaffar, thank you for your support and great effort.

Please find attached sqlquery to extract atm history table to fetch specific day data "Tested in production copy environment (3 Minutes)". Hope this will save more than 5 hours to you. Regards,

Dear All,

We have a good news regarding the OFDM Database Performance.

The Query which used to take 18 minutes is now taking only 27 seconds to get the data from the Database.

Below are the snapshots:

Standard/Golden Rules

Standard/Golden Rules for VLDB configuration

Standard/Golden Rules for VLDB configuration

 Choose a right template in DBCA

-	Database Operation <u>Creation Mode</u>		nplate that include datafiles contain pre-created databases. The n minutes, as opposed to an hour or more. Use templates	
	Database Template	necessary,	such as when you need to change attributes like block si	
-	Database Identification	database c	reation.	
5	Management Options	Select	Template	Includes Datafiles
ļ	Database Credentials	0	General Purpose or Transaction Processing	Yes
	Storage Locations	0	Custom Database Data Warehouse	No Yes
ĺ			testdbs	No
ſ	Database Options			
ľ	Initialization Parameters			
2	Creation Options			
ļ	Pre Réquisite Checks			
	Pre Requisite Checks Summary			
	Summary			
	Summary			
	Summary			Show
	Summary			Show
	Summary			Show

Standard/Golden Rules for VLDB configuration

- ✓ Choose right template
- ✓ Configuring parameters
 - Block Size
 - db_file_multiblock_read_count
 - □ SGA (memory_target/sga_target/pga_aggregate_target)
 - Parallelism
 - Parallel_min_servers
 - Parallel_max_servers
 - Parallel_execution_message_size
 - Parallel_degree_policy
 - o DOP (degree of parallelism)
 - Single instance ,DOP = PARALLEL_THREADS_PER_CPU X CPU_COUNT
 - o RAC , DOP = PARALLEL_THREADS_PER_CPU X CPU_COUNT X INSTANCE_COUNT

Standard/Golden Rules for VLDB configuration

- ✓ Choose right template
- ✓ Configuring parameters
- ✓ Table and Index Partitioning

Partition Strategy Description Range Maps data to partitions based on ranges of partition key values for each partition Hash Maps data to partitions by using a hashing algorithm applied to a partitioning key List Maps data to partitions by using a list of discrete values for the partitioning column Interval Maps data to partitions of ranges that are automatically created by the database following a specified interval System Allows the application to explicitly map rows to arbitrary partitions Composite Is a combination of the basic data-partitioning strategies Index Allows you to choose whether or not to inherit the partitioning strategy of the underlying tables for the indexes Allows the partitioning of two tables related to one another by Reference referential constraints Virtual Column-Based Allows the partitioning key to be defined by an expression, using one or more existing columns of a table

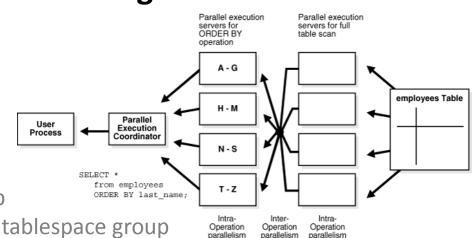
Oracle's Partitioning Strategies

Standard/Golden Rules for VLDB configuration

- ✓ Choose right template
- ✓ Configuring parameters
- ✓ Table and Index Partitioning
- ✓ Avoid or minimize number of indexes
- ✓ Optimal use of index and MV
- ✓ Right index selection, b-tree vs bitmap
- ✓ Bigfile tablespace size and Temporary tablespace groups
- ✓ Data Compression
- ✓ Redo log size and multiplexing
- ✓ Database File System (DBFS)
- ✓ Parallelism

Standard/Golden Rules for VLDB configuration

- ✓ Choose right template
- ✓ Configuring parameters
- ✓ Table and Index Partitioning
- ✓ Avoid or minimize number of indexes
- ✓ Optimal use of index and MV
- ✓ Right index selection, b-tree vs bitmap
- ✓ Bigfile tablespace size and Temporary tablespace group
- ✓ Data Compression
- ✓ Database File System (DBFS)
- ✓ Parallelism
 - Large table scans, joins, or partitioned index will gain performance improvements
 - Large index creation will be benefited
 - Bulk inserts, updates, merges and deletion will gain performance



Optimizer Statistics Management Approach

- ➢ <u>AUTO SAMPLE SIZE</u>
- ➢ INCREMENTAL STATISTICS
- CONCURRENT STATISTICS COLLECTIONS
- LOCK STATS COLLECTION

Optimizer Statistics Management

➢ <u>AUTO SAMPLE SIZE</u>

- ✓ Was much improved in Oracle 11gR1
- ✓ Faster than sampling speed of a 10% sample
- ✓ More accuracy to compute statistics same as 100% accuracy sample
- ✓ Affects histogram gathering and index stats gathering
- ✓ ESTIMATE_PERCENT is set to default AUTO_SAMPLE_SIZE

DBMS_STATS.GATHER_TABLE_STATS(`SCHEMA',`TABLE_NAME',estimate_percent=>DBMS_STATS.AUTO_S
AMPLE_SIZE,granularity=>'AUTO');

Optimizer Statistics Management

Incremental statistics

- Was first introduced in Oracle 11gR1 to improve performance of gathering stats on large partitioned tables
- ✓ Stats for touched partition(s)
- ✓ Global stats are built from partition stats

SQL> select dbms_stats.get_prefs('INCREMENTAL') from dual;

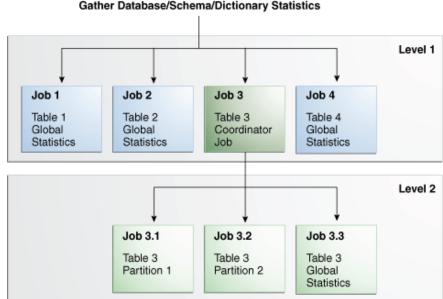
SQL>exec DBMS_STATS.SET_TABLE_PREFS(`SCHEMA',`TABLE_NAME', 'INCREMENTAL','TRUE');

SQL>exec DBMS_STATS.SET_TABLE_PREFS(`SCHEMA',`TABLE_NAME', '**PUBLISH**','TRUE');

Optimizer Statistics Management

Concurrent statistics

- ✓ Was much improved in Oracle 11gR2
- ✓ Gather stats concurrently on multiple tables
- ✓ Goal is to reduce the overall time required for gather stats
- ✓ The global preference CONCURRENT must be set to TRUE
- Oracle uses Job Scheduler and Advanced Queuing (AQ) components to perform the action
- ✓ CREATE JOB, MANAGE SCHEDULER, MANAGE ANY QUEUE
- ✓ JOB_QUEUE_PROCESSES=>4



Optimizer Statistics Management

Concurrent statistics

SQL> ALTER SYSTEM SET RESOURCE_MANAGER_PLAN = 'DEFAULT_PLAN';

SQL> ALTER SYSTEM SET JOB_QUEUE_PROCESSES=8;

SQL> exec BEGIN DBMS_STATS.SET_GLOBAL_PREFS('CONCURRENT', 'ALL MANUAL AUTOMATIC');

SQL> EXEC DBMS_STATS.SET_GLOBAL_PREFS('CONCURRENT', 'OFF');

Backup & Recovery best practices on VLDB

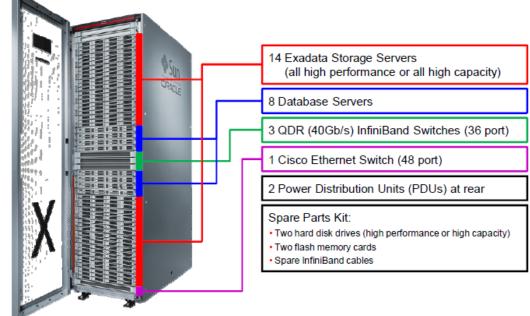
- Developing optimal Backup/Recovery strategies
 - (RTO=recovery time object & RPO=recovery point objective)
- Periodic Full backup and frequent cumulative incremental backups
- Fast incremental backups BCT
- Use parallelism, limit the backup piece size
- Leverage Read only tablespaces and backup compression
- Multi-section backups
- Backup compression
- Create a service for RAC database for backup workload management
- Offload backups to Data Guard
- Other backup solutions, snapshot, split mirror etc
- Compress data
- Backup to FRA (disk) and then to TAPE

Backup approaches

- Level 0 (full) with BCT (Fast Incremental Backups)
- Image copy (full) and Fast Incremental + Incrementally updated backup
- Data Guard + BCT

Exadata solutions

- Hybrid Columnar Compression
- Smart Scan
- Smart Flash Cache
- Smart Flash Cache Log
- IORM



Exadata Key Benefits for Data Warehousing

- Exadata uses more connections.
 - Modular storage cell building blocks are organized into massively parallel grid
 - Bandwidth scales with capacity.
- Exadata has bigger network pipes.
 - InfiniBand network transfers data 5x faster than Fibre Channel.
- Exadata ships less data between storage and database
 - Query processing is moved into storage to dramatically reduce data sent to servers while offloading server CPUs.

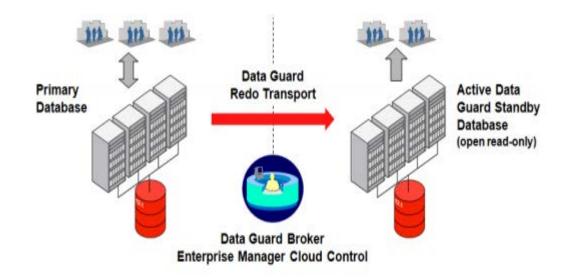
10X to 100X speedup for data warehousing

ASM Best Practices for VLDBs

- Large allocation units
 - Between 1MB and 64MB (4,8,16,32 and 64)
 - Use 64MB for a very large data warehousing system, which will reduce ASM memory requirements, and also improve ASM instance startup time
- Disk groups with an External Redundancy

Active Data Guard Solution

- Offload queries / reports to Standby Database
- Offload backups to Standby Database



Performance optimization

- Indexing (bitmap index)
 - Bitmap indexes are widely used in DW application with huge data and low concurrent DML operations
 - Reduces storage requirements
 - Hardware performance gain, CPU and small amount of memory
- Initialization Parameters
 - OPTIMIZER_MODE = ALL_ROWS | FIRST_ROWS | FIRST_ROWS_n
 - QUERY_REWRITE_ENABLED

A big thank you all for listening ...

You can write me at sjaffarhussain@gmail.com