

Upper Bounds on Geometric Permutations for Convex Sets

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Abstract. Let A be a family of n pairwise disjoint compact convex sets in R^d . Let $\Phi_d(m) = 2\sum_{i=0}^{d-1} \binom{m-1}{i}$. We show that the directed lines in R^d , $d \ge 3$, can be partitioned into $\Phi_d\binom{n}{2}$ sets such that any two directed lines in the same set which intersect any $A' \subseteq A$ generate the same ordering on A'. The directed lines in R^2 can be partitioned into 12n such sets. This bounds the number of geometric permutations on A by $\frac{1}{2}\Phi_d\binom{n}{2}$ for $d \ge 3$ and by 6n for d = 2.

1. Introduction

Let $A = \{a_1, a_2, \ldots, a_n\}$ be a family of n pairwise disjoint compact convex sets in R^d . A common transversal for A is a line intersecting each set in A. Each common transversal intersects the sets in A in a unique order, up to reversal. Katchalski, Lewis, Zaks, and Liu called this ordering and its "reversal" a geometric permutation of A [KLZ], [KLL1]. They showed that any family A of n disjoint compact convex sets in R^2 has at most $\binom{n}{2}$ geometric permutations.

Geometric permutations partition the common transversals of A into disjoint sets such that any two transversals in the same set intersect A in the same order. We can extend the idea of geometric permutations to include orderings from lines which intersect a subset of A. Two lines which intersect $A' \subseteq A$ can be associated if they intersect A' in the same order. We wish to partition the directed lines in R^d into as few sets as possible such that every two lines in the same set which intersect any $A' \subseteq A$ generate the same ordering on A'. Each geometric permutation is associated with a unique pair of sets from the partition. Thus bounding the size of the partition bounds the number of geometric permutations.

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We define a new term, a separation set for A. We say that a hyperplane h separates $a_i \in A$ from $a_j \in A$ if a_i lies in one of the closed half-spaces bounded by h and a_j lies in the other. We say that a hyperplane h strictly separates a convex set a_i from a convex set a_j , if h separates a_i from a_j and $h \cap a_i = h \cap a_j = \emptyset$. If H is a set of hyperplanes in R^d , then we say that H is a separation set for A if, for every pair of elements, a_i , $a_j \in A$, there exists a hyperplane $h \in H$ such that h strictly separates a_i from a_j . Obviously, there exist infinitely many different sets which are separation sets for A.

Let H be a finite set of hyperplanes in \mathbb{R}^d and let |H| be the number of elements in H. Define

$$\Phi_d(m) = 2 \sum_{i=0}^{d-1} {m-1 \choose i}$$

and define $\Psi_d(H)$ to be the number of cones in the pencil obtained by translating the hyperplanes in H to the origin. Winder [W] proved that $\Psi_d(H) \leq \Phi_d(|H|)$.

We present three main theorems in this paper. Bill Lenhart provided the inspiration for Theorem 3.

Theorem 1. Let A be a family of pairwise disjoint compact convex sets in \mathbb{R}^d and H be some separation set for A. The directed lines in \mathbb{R}^d can be partitioned into $\Psi_d(H)$ sets such that every two directed lines in the same set which intersect any $A' \subseteq A$ generate the same ordering on A'.

Theorem 2. If A is a family of n pairwise disjoint convex polygons in \mathbb{R}^2 , then there exists a separation set L for A where each line in L is parallel to some edge of a polygon in A.

Theorem 3. Given A, a family of n pairwise disjoint compact convex sets in \mathbb{R}^2 , there exists a family B of n pairwise disjoint compact convex polygons, such that:

- (i) each convex set in A is entirely contained in a unique polygon in B,
- (ii) the total number of edges in all the polygons in B is at most 12n, and
- (iii) if L is the set of lines containing the edges of the polygons in B, then $\Psi_2(L) \le 12n$.

For every pair of disjoint compact convex sets, there exists some hyperplane which strictly separates the pair [G]. By choosing such a hyperplane for every pair, we can always construct a separation set H for n compact convex sets containing $\binom{n}{2}$ elements. Since $\Psi_d(H) \leq \Phi_d(|H|)$, we have the following corollary to Theorem 1:

Corollary 1. If A is a family of n pairwise disjoint compact convex sets in \mathbb{R}^d , the directed lines in \mathbb{R}^d can be partitioned into $\Phi_d \binom{n}{2}$ sets such that any two directed lines in the same set which intersect any $A' \subseteq A$ generate the same ordering on A'.

Finally, combining Theorems 2 and 3, we can always find a separation set L for n convex sets in R^2 such that $\Psi_2(L) \le 12n$. Applying Theorem 1, we have the following corollary:

Corollary 2. If A is a family of n pairwise disjoint compact convex sets in \mathbb{R}^2 , the directed lines in \mathbb{R}^2 can be partitioned into 12n sets such that any two directed lines in the same set which intersect any $A' \subseteq A$ generate the same ordering on A'.

As noted before, each geometric permutation is associated with a unique pair of sets in the partition of directed lines. Thus the number of geometric permutations on A is bounded by half the size of the partition. Corollaries 1 and 2 imply that there are at most $\frac{1}{2}\Phi_d\binom{n}{2}$ geometric permutations of A in R^d and at most 6n geometric permutations of A in R^d .

2. Upper Bounds on the Number of Geometric Permutations

We now prove Theorem 1. Let $A = \{a_1, a_2, \ldots, a_n\}$ be a family of n pairwise disjoint compact convex sets in R^d and let $H = \{h_1, h_2, \ldots, h_m\}$ be a separation set of A containing m elements. Let $\{u_1, u_2, \ldots, u_m\}$ be a set of normal vectors to H where u_k is normal to h_k .

Let v be a vector in R^d and let T be the set of all lines with direction v. Let h_k be the hyperplane separating a_i and a_j and assume $(a_j - a_i) \cdot u_k > 0$. If $v \cdot u_k > 0$, then any line in T must intersect a_i before a_j . If $v \cdot u_k < 0$, then any line in T must intersect a_j before a_i . Furthermore, even if no line with direction v intersects a_i and a_j simultaneously, $\operatorname{sgn}(v \cdot u_k)$ still determines an ordering relation on a_i and a_j . Thus $\operatorname{sgn}(v \cdot u_k)$, $k = 1, \ldots, m$, determines the relative order in which any line with direction v intersects any two elements of A.

Consider the cones created by hyperplanes through the origin with normals u_k , $k=1,\ldots,m$. The values of $\operatorname{sgn}(v\cdot u_k)$, $k=1,\ldots,m$, are determined by the cone into which v points. Any two lines which point into a given cone and intersect some $A'\subseteq A$ generate the same ordering on A'. There are at most $\Psi_d(H)$ cones, and so the directed lines can be partitioned into $\Psi_d(H)$ sets with the desired property. This proves Theorem 1.

3. Separation Sets for Convex Polygons

Let A be a family of pairwise disjoint compact convex polygons in R^2 . For Theorem 2, we show how to construct a separation set L for A where each line in L is parallel to some edge in A. We need merely show that for any two polygons, a_i and a_j , there exists some line strictly separating a_i from a_j and parallel to some edge in A.

Let l^* be some line strictly separating convex polygon a_i from convex polygon a_j and tangent to a_i at some vertex. Let l_1 and l_2 be the two lines containing the

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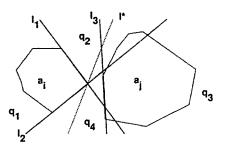


Fig. 1. Convex polygons a_i and a_j .

edges of a_i which meet at that vertex. l_1 and l_2 divide the plane into four quadrants, q_1 , q_2 , q_3 , q_4 , with a_i lying wholly in one quadrant, say q_1 (see Fig. 1). We include in each quadrant the boundary of that quadrant.

Since l^* separates a_j from a_i , a_j does not intersect quadrant q_1 . If a_j does not intersect quadrant q_2 , then l_2 separates a_i from a_j . If a_j does not intersect quadrant q_4 , then l_1 separates a_i from a_j . By translating either of these separators slightly toward a_j , we have a line strictly separating a_i from a_j and parallel to a line through an edge of a_i .

Assume a_j contains a point from quadrant q_2 and from quadrant q_4 . Some edge of a_j must intersect l_1 . Let l_3 be the line containing that edge. l_3 separates a_1 from a_2 so by translating it slightly toward a_j we have a line strictly separating a_i from a_j and parallel to an edge of a_i .

For each pair of polygons a_i , $a_j \in A$, we add to L a line strictly separating a_i from a_j and parallel to some edge in a_i or a_j . L is a separation set for A such that every line in L is parallel to some edge of a polygon in A, fulfilling the conditions in Theorem 2.

4. Embedding Compact Convex Sets in Polygons

Given n pairwise disjoint compact convex sets, we show how to embed these convex sets in n pairwise disjoint compact convex polygons satisfying the conditions in Theorem 3. We first note that n pairwise disjoint convex sets can be embedded in n pairwise disjoint convex polygons using a total of not more than n(n+2) edges. Choose some convex set a and find n-1 lines separating a from the n-1 other convex sets. These n-1 lines bound n-1 half-planes containing a. If the intersection of these half-planes is unbounded, we can add three suitable half-planes containing a such that the intersection of the n+2 half-planes is bounded. The intersection of these half-planes forms a polygon with at most n+2 edges which contains a. Repeat the procedure n times to get a set a of a pairwise disjoint convex polygons containing the a convex sets.

Following [PS], a complete triangulation of a finite set of points S is a planar graph whose vertices are the points of S, whose edges are nonintersecting line segments joining the points in S, and whose faces (including the external face) are bounded by triangles. We define a new type of triangulation, a complete



Fig. 2. A complete triangulation of convex polygons.

triangulation of a finite family of convex polygons A is a planar graph whose vertices can be mapped to polygons in A, whose edges can be mapped to nonintersecting line segments joining two polygons and not intersecting any other polygon, and whose faces (including the external face) are bounded by triangles (see Fig. 2). We allow graphs to have multiple edges, i.e., more than one edge between any two vertices.

Let c be some polygon in A which lies on the convex hull of A. Add two small triangles, d and d', to A such that the convex hull of $A' = A \cup \{d, d'\}$ is three line segments between c, d, and d'. The convex hull of A' minus the polygons in A' forms a polygon with holes. Triangulate this polygon with holes [PS]. This triangulation induces a planar graph on the polygons in A' where polygons are mapped to vertices and triangulation line segments to edges. The faces in this planar graph are bounded by two or three edges. If a face is bounded by only two edges, one of the bounding edges and its corresponding line segment is redundant. Remove these redundant edges and line segments until all faces are bounded by three edges. The resultant planar graph G is a complete triangulation of A'.

Given this triangulation G of A' we can embed the n pairwise disjoint convex polygons in A into n pairwise disjoint convex polygons with at most 12n edges. For each triangulation line segment s connecting a, $b \in A'$, let l_s be the line containing s. Let $l_{a,b}$ be some line strictly separating a from b.

For each convex set $a \in A'$, let N(a) be a list of the neighbors of a. Let F(a) be a list of the faces which have a as a vertex. Let $L_1(a) = \{l_{a,b} : b \in N(a)\}$ and let $L_2(a) = \{l_s : s \text{ lies on a face of } F(a) \text{ and connects } b, b' \in N(a)\}$. Let $L_0(a)$ be the union of the lines in $L_1(a)$ and $L_2(a)$ slightly translated toward a. Let P(a) be the intersection of the half-planes containing a and bounded by the lines in $L_0(a)$. For any $a \in A$ other than c, P(a) is bounded by the cycle of the neighbors of a and the triangulation line segments between them (see Fig. 3). Thus P(a) is a convex polygon containing a. By a judicious choice of $l_{c,d}$ and $l_{c,d'}$, P(c) is also a convex polygon containing c. We claim that $B = \{P(a) : a \in A\}$ is a family of n pairwise disjoint compact convex polygons such that each $a \in A$ is contained in P(a) and the total number of edges in the polygons in B is 12n.

Assume $a, b \in A$ are neighbors in G. P(a) and P(b) are both bounded by a translation of the same separator, $l_{a,b}$ of a and b. Since P(a) and P(b) are in different half-planes bounded by $l_{a,b}$, $l_{a,b}$ must separate P(a) from P(b) and P(a) and P(b) are pairwise disjoint.

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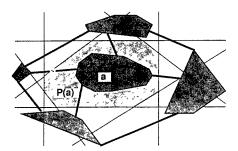


Fig. 3. Cycle of neighbors of a.

Now, assume $a, b \in A$ are not neighbors in G. P(a) is entirely contained in a cycle of the neighbors of a while P(b) is entirely contained in a cycle of the neighbors of b. Since b is not a neighbor of a, these two cycles must contain different regions in the plane and so P(a) and P(b) must be pairwise disjoint.

Let m be the number of edges in G. G is a planar graph in which every face is a triangle and there are n+2 vertices in G, so m=3n [H]. Each edge connects two neighbors, so the total number of neighbors over all vertices in the graph is 6n. The total number of edges in the polygons in B is at most the total number of lines in $L_0(a)$ over all $a \in A$:

$$\sum_{a \in A} |L_0(a)| \le \sum_{a \in A'} |L_0(a)| = \sum_{a \in A} |L_1(a)| + \sum_{a \in A} |L_2(a)| \le 6n + 6n = 12n.$$

Therefore, the total number of edges in the polygons in B is at most 12n.

Let L be the set of lines containing the edges of the polygons in B. L is a subset of $L' = \bigcup_{a \in A'} L_0(a)$. Each line in $L_0(a)$ is parallel to some line in some $L_0(b)$, $a, b \in A'$. Since the total number of lines in L' is at most 12n, $\Psi_2(L') \le 12n$ and so $\Psi_2(L) \le 12n$.

5. Conclusion

Katchalski et al. [KLZ] asserted that, for every d, there exists a constant k_d and a family A of n pairwise disjoint compact convex sets in R^d such that there are at least $k_d n^{d-1}$ geometric permutations of A. Corollary 1 implies an upper bound of $O(n^{2d-2})$ for the number of geometric permutations of A, leaving a wide gap for improvement. Villanger showed that for any n there exist families of n line segments in R^3 in which any hyperplane separates at most one line segment from one other line segment [T]. Thus, there exist a family A of n compact convex sets in R^d , for any $d \ge 3$, such that any separation set of A must have $\binom{n}{2}$ elements. In fact, by embedding Villanger's line segments in rectangular prisms, we see that Theorem 2 does not generalize to two dimensions. Reduction of the upper bounds must come from other directions.

Much attention has been paid to geometric permutations of families of translates [K], [KLL1], [KLL2]. Are there separation theorems for families of translates? These theorems may increase our understanding about geometric permutations of families of translates.

We are also interested in reducing the 12n edges needed for n pairwise disjoint compact convex polygons to contain n pairwise disjoint compact convex sets in R^2 . Approximately 6n edges are needed for any n pairwise disjoint convex polygons containing n hexagons in a lattice of hexagons. We do not know if any set of n compact convex sets in R^2 can be contained in n compact convex polygons in R^2 with a total of about 6n edges. Finally, we do not know if Theorem 3 generalizes to three dimensions. Can any set of n pairwise disjoint compact convex sets in R^3 be embedded in n pairwise disjoint convex polyhedra such that the total number of faces is linear in n?

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